

D_{RS} -TRIANGULATION BASED ON 2^N -RAY ALGORITHM

By

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*(Received : November 05, 2010; Revised: March 20, 2011)***ABSTRACT**

In order to improve the efficiency of the 2^N -Ray algorithm; we propose a variant of the D_{RS} -Triangulation. A nice property of this triangulation is that it subdivides all the subsets, on which the 2^N -Ray algorithm works, into simplices according to the D_{RS} -Triangulation. Numerical tests shows that 2^N -Ray algorithm based on $D_{1/2}$ -Triangulation is much more efficient.

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1. Introduction. Various triangulations have been studied by Dang [1], Dang and Talman [2], Scarf [4], Todd [6,7], Vander and Talman [8], Vertgeim [9]. The 2^n -ray algorithm was proposed by Wright in [10] to compute solutions of non-linear equations. The 2^n -ray algorithm partitions R^n into $2n$ cones which have the same vertex. Then a triangulation of R^n subdivides each cone into simplices. The 2^n -ray algorithm starts at the vertex and leaves it along an edge of some cone. It follows a sequence of adjacent simplices with varying dimension. Under some mild conditions, the 2^n -ray algorithm terminates at an n -dimensional simplex that yields an approximate solutions to the system of non-linear equations. Since these $2n$ cones have 2^n edges, each of which is a ray, the 2^n -ray algorithm has 2^n possible ray to leave the vertex.

Motivated by this work, we have tried to develop a new D_{RS} -triangulation. The D_{RS} -triangulation based on the 2^n ray algorithm is much more efficient than 2^n ray algorithm proposed by Wright [10] as well as K_1 triangulation introduced by Kuhn in [3] and J_1 triangulation given by Todd [5].

2. Notations and Definitions. The following notations have been used in this paper :

R : Set of real numbers,

Z : Set of all integers,

N : Set of positive integers $(1,2,\dots,n)$,

N^0 : Set of all non-negative integers $N \cup (0)$,

- R^n : n dimensionally space, having co-ordinates indexed 1 through n ,
 R^{n+1} : $n+1$ dimensional space, with coordinates indexed 0 through n ,
 π : Group of permutation on $(1,2,\dots,n)$ and $\pi+1$ group of permutation on $(0,1,2,\dots,n)$,
 u^i : i^{th} unit vector in R^n , $j \in N$ and $u = \sum_{i \in N} u^i$,
 R_+^m : Non negative orthant of R^m i.e. $(x \in R^m; x \geq 0)$.

Now we consider some standard definitions and explanations which will be used in this paper.

2.1 Standard Simplex. The standard n dimensional closed simplex S^n is the convex hull of v^0, v^1, \dots, v^n i.e. $S^n = \{x \in R_+^{n+1} : v^T x = 1\}$. s_i^n denotes the face of s^n opposite v^i i.e. $S_i^n = \{x \in s^n : x_i = 0\}$ and boundary of s^n is denoted by $\partial s = \bigcup_{i \in N} s_i^n$.

Again a j -dimensional simplex or [j -simplex] is the relative interior of the convex hull of $j+1$ affinely independent points $y^0, y^1, y^2, y^3, \dots, y^j$, called its vertices.

We write $\sigma = \langle y^0, y^1, y^2, \dots, y^j \rangle$. A simplex τ is a face of σ if its vertices are a subset vertices of the σ . It is convenient to call the closure of a $(j-1)$ dimensional face of the j simplex σ as a facet of σ . Two j simplices are said to be adjacent if they share a common facet.

2.2 Triangulation. A triangulation G of S^n is a collection of n simplices and satisfies the following two conditions :

1. The simplices in G together with all their faces form a partition of S^n and
2. Each point of S^n has a neighbourhood meeting only a finite number of simplices.

2(a) Pivot Rule. For a given simplices G and a vertex y of σ the rules for obtaining the simplex of G whose vertices include all vertices of σ except y , are called the pivot rules of G .

2(b) Mesh. The mesh of a triangulation G is $\sup_{\sigma \in G} \text{diam} \sigma$. We shall use the Euclidian norm through out this paper.

2.3 Definition. For each sign vector $s \in R^n$, let

$$\begin{aligned}
 E(s) &= \{x \in R^n : s_i x_i = \|x\| \text{ whenever } s_i \neq 0\} \\
 &= \text{cone} \{t \in R^n : t \text{ is a sign vector, and } s_i \neq 0 \Rightarrow s_i = t_i\} .
 \end{aligned}$$

In case s has k non-zero components for $k > 0$ than $E(s)$ is a polyhedral cone of dimension $n - k + 1$. Also we have $E(0) = R^n$. Moreover, when $s \neq 0$ each $E(s) \cap B^n$

is a polyhedral of a cubical subdivisions of B^{n_∞} where B^{n_∞} denote the unit ball in 1^∞ norm.

Wright [10] has also defined another subdivision of R^n into closed convex cone as n -dimensional geometric form given as follows :

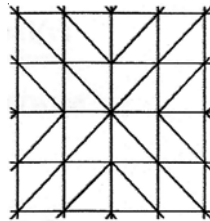
$$\text{Let } C(s) = x \in R^n : \begin{cases} x_i = 0 \text{ if } s_i = 0 \\ s_i x_i \geq 0 \text{ if } s_i \neq 0 \end{cases}$$

= cone $\{s_i u_i : s_i > 0, \text{ for each sign vector } s\}$. If s has k non-zero components then $C(s)$ is an orthant of a k -dimensional coordinate subspace of R^n .

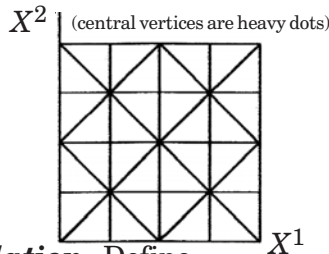
Wright [10] has proposed two type of T -triangulations of R^n with the property that $E(t)$ is a subcomplex for every sign vector $t \in R^n$ for $t \neq 0$. The first triangulation is called a K^1 triangulation. This triangulation is obtained by taking the triangulation K_1 due to Kuhn [3] in the first orthant and the reflecting through coordinate hyperplane to triangulate the other orthant. A vector v^1 of a n -simplex $\langle v^1, \dots, v^{n+1} \rangle$ of K^1 specified by choosing a sign vector s with all the non-zero components, is a member of $C(s)$ when all its components are integrals and π is a permutation of $\{1, 2, \dots, n\}$, then $v^1 \in K^1$ is defined recursively as :

$$V^{i+1} = V^i + S_{\pi(i)} U^{\pi(i)}, \text{ for } i=1, 2, \dots, n.$$

For $n=2$ triangulation K^1 can be illustrated by following diagram.



The second triangulation J_1 is defined by Todd [5] and which can be illustrated for $n=2$ by the following figure :



2.3.1. The D_{RS} -Triangulation. Define

$$W^n = \{x \in R_+^n : x_1 = \max x_i, i = 2, 3, \dots, n\}$$

taking a vector $Y = (y_1, y_2, \dots, y_n)^T$, we have

$$Y_i = \begin{cases} \lceil x_i \rceil & \text{if } \lceil x_i \rceil \text{ is even} \\ \lceil x_i \rceil + 1 & \text{otherwise,} \end{cases}$$

where $\lceil \alpha \rceil$ is the greatest integer less than or equal to α . Let D is the set of all $Y \in W^n$ where Y_i is defined above. If $Y \in D$, we define

$$I(y) = \{i \in N; y_i = y_i\} \text{ and } J(y) = \{j \in N : y_1 \geq y_j\}^T.$$

Let $s = (s_1, s_2, \dots, s_n)^T$ be a sign vector such that

1. For $i \in N$, if $y_i = 0$ then $s_i = 1$, and if $y_i \neq 0$ then $s_i = -1$.

$$\text{Let } K(y, s) = \{i \in I(y) : s_i = 1\}.$$

Let ℓ denote the number of element in $I(y)$ and h the number of elements in $K(y, s)$, we take integer p such that

1. when $h=0$, if $\ell = n$ then $p=0$ or 2 ,
2. when $h>0$, if $h=n$ then $p=0$ and if $h<n$ then $0 \leq p \leq n-1$.

Let $\pi = \{\pi(1), \pi(2), \dots, \pi(n)\}$ be permutation of N .

When $h=0$, for $j=1, 2, \dots, n$,

1. If $j=1$, define

$$g_i(j) = \begin{cases} -1 & \text{if } i \in I(y) \\ 0 & \text{otherwise} \end{cases} \quad \dots(1)$$

for $i=1, 2, \dots, n$.

2. If $j \neq 1$, we define

$$g_i(j) = \begin{cases} S^i & \text{if } i = j, \\ 0 & \text{otherwise} \end{cases} \quad \dots(2)$$

for $i=1, 2, \dots, n$.

When $h>0$, for $j=1, 2, \dots, n$.

1. If $\pi(j) \in K(y, s)$, define

$$g_i(\pi(j)) = \begin{cases} 1 & \text{if } i \in K(y, s) \text{ and } j \leq \pi^{-1}(i) \\ 0 & \text{otherwise} \end{cases} \quad \dots(3)$$

for $i=1, 2, \dots, n$,

2. If $\pi(j) \notin K(y, s)$, define

$$g_i(\pi(j)) = \begin{cases} s_n(j) & \text{if } i = \pi(j) \\ 0 & \text{otherwise} \end{cases} \quad \dots(4)$$

for $i=1,2,\dots,n$.

If y, π, s and p be as above, then vectors y^0, y^1, \dots, y^n are defined as follows :

for $p=0$, we have

$$\begin{aligned} y^0 &= y \\ y^k &= y + g(\pi(k)), k = 1, 2, \dots, n \end{aligned} \quad \dots(5)$$

and for $p \geq 1$, we define

$$\begin{aligned} y^0 &= y + s, \\ y^k &= y^{k-1} - s_{\pi(k)} U^{(\pi(k))}, k = 1, 2, \dots, p-1. \\ y^k &= y + g(\pi(k)), k = p, \dots, n. \end{aligned} \quad \dots(6)$$

The y^0, y^1, \dots, y^n vectors obtained from the above definition are affinely independent.

Thus their convex hull is a simplex. Let us denote this simplex by $D_{RS}(y, \pi, s, p)$ or

$\langle y^0, y^1, \dots, y^n \rangle$. Let D_{RS} be the set of all such simplices. Then D_{RS} is a triangulation of W^n . Note that simplices of the D_{RS} -triangulation can be represented in more than one way. Moreover, the triangulation of a whole cube in W^n is the same as the D_1 -triangulation.

To be more expatiate let us illustrate D_{RS} -triangulation of W^n for $n=2$ and for $x \leq 4$. Obviously, we have that for $y_1 \leq 4$.

$$D = \left\{ (0,0,0)^T, (2,2,0)^T, (4,4,0)^T, (4,0,4)^T, (4,4,4)^T \right\}. \quad \dots(7)$$

1. Let $y = (0,0,0)^T$. Then, $I(y) = \{i \in N : y_i = y_i\} = (1,2,3)$ and $\ell = 3$.

Then s must be $(1,1,1)^T$. Thus

$K(y, s) = \{i \in I(y) : s_i = 1\} = (1,2,3)$ and $h = 3$. We have $p = 0$.

(a) Let $\pi = (2,3,1)$. Then $\pi^{-1} = (3,1,2)$ and by applying (3)

$$g(\pi(1)) = g(2) = (g_1(2), g_2(2), g_3(2))$$

$$= (1,1,1)^T,$$

$$g(\pi(2)) = g(3) = (1,0,1)^T,$$

$$g(\pi(3)) = g(3) = (1,0,0)^T.$$

Therefore, $y^0 = y = (0,0,0)^T,$

$$y^1 = y + g(\pi(1)) = (1,1,1)^T,$$

$$y^2 = y + g(\pi(2)) = (1,0,1)^T,$$

$$y^3 = y + g(\pi(3)) = (1,0,0)^T, \text{ Let } \sigma^1 = \langle y^0, y^1, y^2, y^3 \rangle.$$

(b) Let $\pi = (3,2,1)$. Then $\pi^{-1} = (3,2,1)$

$$g(\pi(1)) = g(3) = (1,1,1)^T,$$

$$g(\pi(2)) = g(2) = (1,1,0)^T,$$

$$g(\pi(3)) = g(1) = (1,0,0)^T,$$

Therefore, $y^0 = y = (0,0,0)^T,$

$$y^1 = y + g(\pi(1)) = (1,1,1)^T,$$

$$y^2 = y + g(\pi(2)) = (1,1,0)^T,$$

$$y^3 = y + g(\pi(3)) = (1,0,0)^T, \text{ Let } \sigma^2 = \langle y^0, y^1, y^2, y^3 \rangle.$$

2. Let $y = (2,2,0)^T$. Since $I(y) = \{i \in N : y_1 = y_i\} = \{1,2\}$ and $\ell = 2$.

So s must be $(-1,-1,1)^T$. Thus $k(y,s) = \{i \in I(y); s_i = 1\} = \emptyset$ and $h=0$ while p can be any one of $0,1,2$.

Now by applying (1) and (2), we have

$$g(1) = (-1,-1,0)$$

$$g(2) = (0,-1,0)$$

$$g(3) = (0,0,1).$$

Now considering different values of p and applying (6) and (7), we obtain the following simplices:

(a) For $p=0$. Let $\pi = (1,2,3)$. Therefore,

$$\begin{aligned} y^0 &= y = (2,2,0)^T, \\ y^1 &= y + g(\pi(1)) = (1,1,0)^T, \\ y^2 &= y + g(\pi(2)) = (2,1,0)^T, \\ y^3 &= y + g(\pi(3)) = (2,2,1)^T. \text{ Let } \sigma^3 = \langle y^0, y^1, y^2, y^3 \rangle. \end{aligned}$$

(b) For $p=1$. Let $\pi = (1,2,3)$. Therefore

$$\begin{aligned} y^0 &= y + s = (1,1,1)^T, \\ y^1 &= y + g(\pi(1)) = (1,1,0)^T, \\ y^2 &= y + g(\pi(2)) = (2,1,0)^T, \\ y^3 &= y + g(\pi(3)) = (2,2,1)^T. \text{ Let } \sigma^4 = \langle y^0, y^1, y^2, y^3 \rangle. \end{aligned}$$

(c) For $p=2$. Let $\pi = (1,2,3)$. We have,

$$\begin{aligned} y^0 &= y + s = (1,1,1)^T, \\ y^1 &= y^0 - s_{\pi(1)} u^{\pi(1)} = (2,1,1)^T, \\ y^2 &= y + g(\pi(2)) = (2,1,0)^T, \\ y^3 &= y + g(\pi(3)) = (2,2,1)^T. \text{ Let } \sigma^5 = \langle y^0, y^1, y^2, y^3 \rangle. \end{aligned}$$

3. Let $y = (4,4,0)^T$. Therefore, $I(y) = \{i \in N; y_i = y_i\} = \{1,2\}$ and $\ell = 2$.

We have that s must be $(-1,-1,1)^T$. Thus $K(y,s) = \{i \in I(y) : s_i = 1\} = \emptyset$ and $h=0$.

We have that p can be any one of 0,1,2. We also have

$$\begin{aligned} g(1) &= (-1,-1,0), \\ g(2) &= (0,-1,0), \\ g(3) &= (0,0,1). \end{aligned}$$

(a) For $p=0$, and $\pi = (1,2,3)$, we have

$$\begin{aligned} y^0 &= y = (4,4,0)^T, \\ y^1 &= y + g(\pi(1)) = (3,3,0)^T, \\ y^2 &= y + g(\pi(2)) = (4,3,0)^T, \\ y^3 &= y + g(\pi(3)) = (4,4,1)^T. \text{ Let } \sigma^6 = \langle y^0, y^1, y^2, y^3 \rangle. \end{aligned}$$

(b) For $p=1$, let $\pi = (1,2,3)$. Therefore

$$y^0 = y + s = (3,3,1)^T,$$

$$y^1 = y + g(\pi(1)) = (3,3,0)^T,$$

$$y^2 = y + g(\pi(2)) = (4,3,0)^T,$$

$$y^3 = y + g(\pi(3)) = (4,4,1)^T. \text{ Let } \sigma^7 = \langle y^0, y^1, y^2, y^3 \rangle.$$

(c) For $p=2$ and $\pi = (1,2,3)$, we have

$$y^0 = y + s = (3,3,1)^T,$$

$$y^1 = y^0 - s_{\pi(1)} u^{\pi(1)} = (4,4,1)^T,$$

$$y^2 = y + g(\pi(2)) = (4,3,0)^T,$$

$$y^3 = y + g(\pi(3)) = (4,4,1)^T. \text{ Let } \sigma^8 = \langle y^0, y^1, y^2, y^3 \rangle.$$

4. Let $y = (4,0,4)^T$. Since $I(y) = \{i \in N; y_i = 1\} = \{1,3\}$ and $\ell = 2$. so that s must be $(-1,-1,1)^T$. Thus $K(y,s) = \{i \in I(y) : s_i = 1\} = \emptyset$ and $h=0$. we have that p can be any one of $0,1,2$. We also have

$$g(1) = (-1,0,-1),$$

$$g(2) = (0,1,0),$$

$$g(3) = (0,0,-1).$$

(a) For $p=0$, and $\pi = (1,2,3)$, we have

$$y^0 = y = (4,0,4)^T,$$

$$y^1 = y + g(\pi(1)) = (3,0,3)^T,$$

$$y^2 = y + g(\pi(2)) = (4,1,4)^T,$$

$$y^3 = y + g(\pi(3)) = (4,0,3)^T. \text{ Let } \sigma^9 = \langle y^0, y^1, y^2, y^3 \rangle$$

(b) For $p=1$. Let $\pi = (1,2,3)$. Therefore

$$y^0 = y + s = (3,1,3)^T,$$

$$y^1 = y + g(\pi(1)) = (3,0,3)^T,$$

$$y^2 = y + g(\pi(2)) = (4,1,4)^T,$$

$$y^3 = y + g(\pi(3)) = (4,0,3)^T. \text{ Let } \sigma^{10} = \langle y^0, y^1, y^2, y^3 \rangle$$

(c) For $p=2$ and $\pi = (1,2,3)$, we have

$$\begin{aligned}
y^0 &= y + s = (3, 1, 3)^T, \\
y^1 &= y^0 - s_{\pi(1)} u^{\pi(1)} = (4, 1, 3)^T, \\
y^2 &= y + g(\pi(2)) = (4, 1, 4)^T, \\
y^3 &= y + g(\pi(3)) = (4, 0, 3)^T. \text{ Let } \sigma^{11} = \langle y^0, y^1, y^2, y^3 \rangle
\end{aligned}$$

5. Let $y = (4, 4, 4)^T$. Since $I(y) = \{i \in N; y_i = y_i\} = \{1, 2, 3\}$ and $\ell = 3$. We have that a must be $(-1, -1, -1)^T$. Thus $K(y, s) = \emptyset$ and $h = 0$. we have that p can be any one of 0, 1, 2. We also have

$$\begin{aligned}
g(1) &= (-1, -1, -1), \\
g(2) &= (0, -1, 0), \\
g(3) &= (0, 0, -1).
\end{aligned}$$

(a) For $p=0$. Let $\pi = (1, 2, 3)$. Therefore

$$\begin{aligned}
y^0 &= y = (4, 4, 4)^T, \\
y^1 &= y + g(\pi(1)) = (3, 3, 3)^T, \\
y^2 &= y + g(\pi(2)) = (4, 3, 4)^T, \\
y^3 &= y + g(\pi(3)) = (4, 4, 3)^T. \text{ Let } \sigma^{12} = \langle y^0, y^1, y^2, y^3 \rangle.
\end{aligned}$$

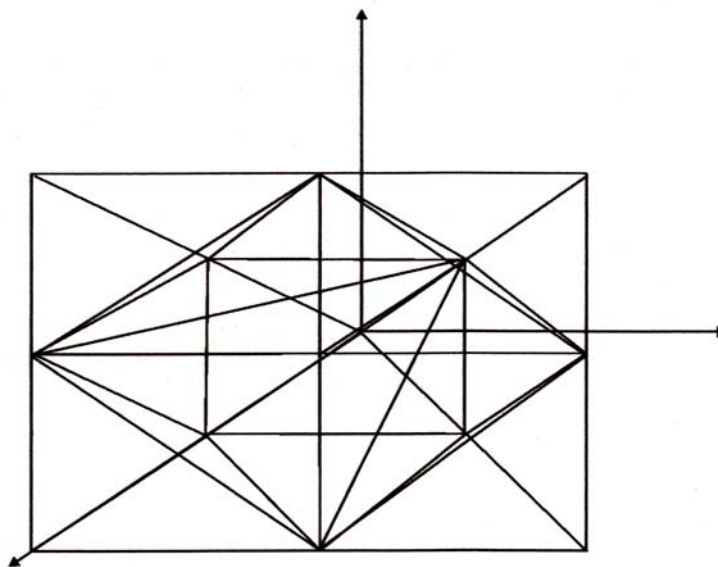
(b) For $p=1$. Let $\pi = (1, 2, 3)$. Therefore

$$\begin{aligned}
y^0 &= y + s = (3, 3, 3)^T, \\
y^1 &= y + g(\pi(1)) = (3, 3, 3)^T, \\
y^2 &= y + g(\pi(2)) = (4, 3, 4)^T, \\
y^3 &= y + g(\pi(3)) = (4, 4, 3)^T. \text{ Let } \sigma^{13} = \langle y^0, y^1, y^2, y^3 \rangle.
\end{aligned}$$

(c) For $p=2$. Let $\pi = (1, 2, 3)$. We have,

$$\begin{aligned}
y^0 &= y + s = (3, 3, 3)^T, \\
y^1 &= y^0 - s_{\pi(1)} u^{\pi(1)} = (4, 3, 3)^T, \\
y^2 &= y + g(\pi(2)) = (4, 3, 4)^T, \\
y^3 &= y + g(\pi(3)) = (4, 4, 3)^T. \text{ Let } \sigma^{14} = \langle y^0, y^1, y^2, y^3 \rangle
\end{aligned}$$

It can be seen that $\{\sigma^i : i = 1, 2, \dots, 14\}$ form a triangulation of W^3 for $x \leq 4$ as shown by the following figure :



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